

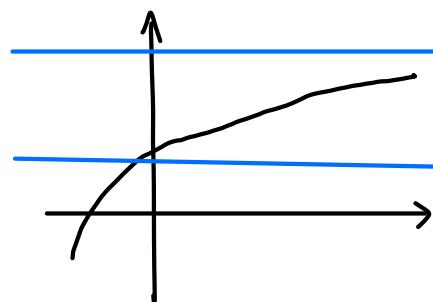
Lastnosti preslikav

$$f: A \rightarrow B$$

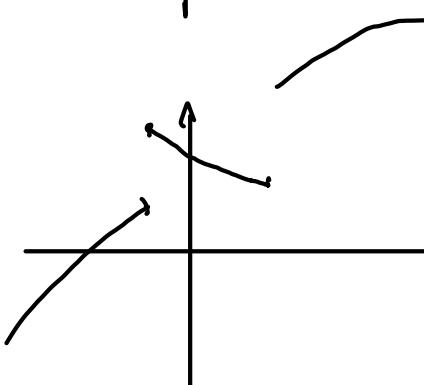
• injektivna : $\forall x, y \in A . f(x) = f(y) \Rightarrow x = y$

ekvivalentno:

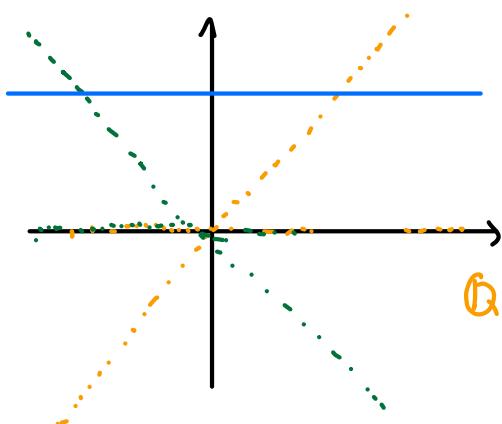
$$\forall x, y \in A . x \neq y \Rightarrow f(x) \neq f(y)$$



seka graf
največ enkrat



injektivna
(ni naravljajoča,
ni padajoča)

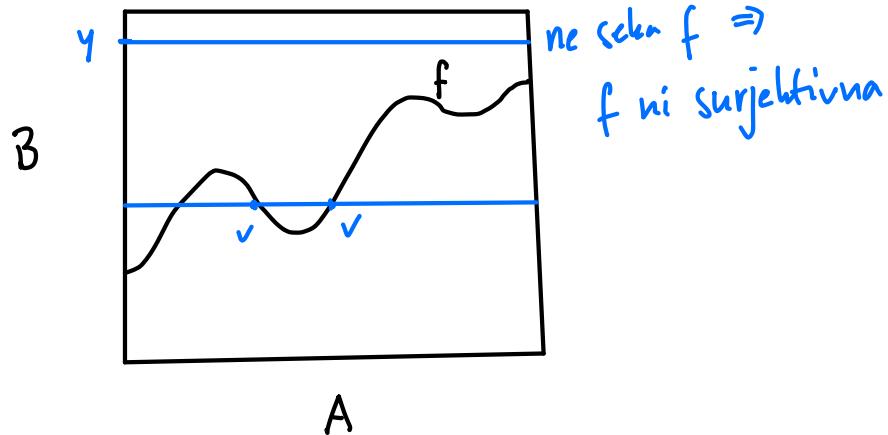


$$f: \mathbb{R} \rightarrow \mathbb{R}$$

$$f(x) = \begin{cases} x & \text{če je } x \in \mathbb{Q} \\ -x & \text{če je } x \notin \mathbb{Q} \end{cases}$$

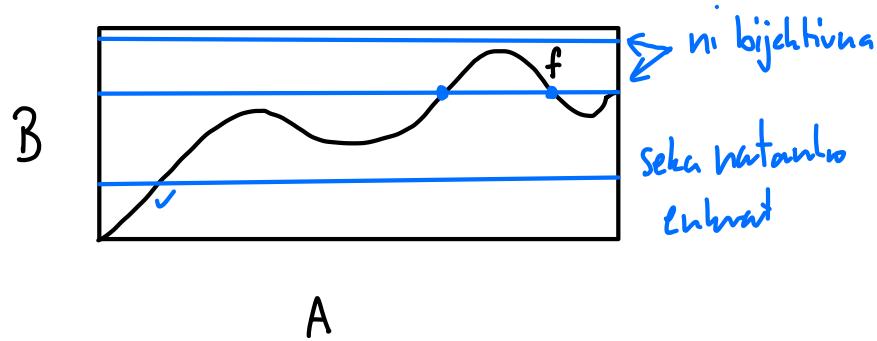
$$f: A \rightarrow B$$

Surjektivna: $\forall y \in B. \exists x \in A. y = f(x)$.



bijektivna: injektivna & surjektivna

ekvivalenten: $\forall y \in B. \exists ! x \in A. y = f(x)$



Monomorfizmi & epimorfizmi

matrice

$$A \cdot B = A \cdot C ? \text{ ali smemo kerajšati?}$$
$$B = C$$

$$B \cdot A = C \cdot A ?$$

Def : Preslikava $f: A \rightarrow B$ je

- monomorfizem (mono), ko jo lahko krajšamo na levi:

$$\forall C \in \text{Set}. \underbrace{\forall g, h \in A^C}_{\text{zapis}}. f \circ g = f \circ h \Rightarrow g = h$$

zapis : $\forall g, h: C \rightarrow A.$

$$C \xrightarrow{\begin{matrix} g \\ h \end{matrix}} A \xrightarrow{f} B$$

$$\forall \varepsilon > 0. \varphi \text{ dvojčaka za}$$

$$\forall \varepsilon \in \mathbb{R}. \varepsilon > 0 \Rightarrow \varphi$$

$$\exists \delta > 0. \varphi \text{ za}$$

$$\exists \delta \in \mathbb{R}. \delta > 0 \wedge \varphi$$

- epimorfizem (epi), ko jo lahko krajšamo na desni

$$\forall D \in \text{Set}. \forall g, h: B \rightarrow D. g \circ f = h \circ f \Rightarrow g = h$$

SMEŠNO

$$A \xrightarrow{f} B \xrightarrow{\begin{matrix} g \\ h \end{matrix}} D$$

Izrek ⑦ Naj bosta $A \xrightarrow{f} B \xrightarrow{g} C$. Velja:

1) Kompozicija monomorfizmov je monomorfizem

2) $\dashv \dashv$ epi $\dashv \dashv$ epi

3) Če je $g \circ f$ mono, potem je f mono.

4) Če je $g \circ f$ epi, potem je g epi.

Dokaz.

1) Predpostavimo f mono
 g mono

$$C \xrightarrow{k} A \xrightarrow{f} B \xrightarrow{g} C$$

$\underbrace{\hspace{1cm}}_{g \circ f}$

Dohazimo $g \circ f$ mono:

Naj bu $C \in \text{Set}$.

Naj boste $h, k: C \rightarrow A$.

Dohazimo: $(g \circ f) \circ h = (g \circ f) \circ k \Rightarrow h = k$

Predpostavimo: $(g \circ f) \circ h = (g \circ f) \circ k$
(asociativnost) $\quad \parallel \quad \parallel$

$$\begin{array}{ccc} (g \circ f) \circ (f \circ h) & = & g \circ (f \circ h) \\ (g \text{ mono}) \quad \cancel{f \circ h} & = & f \circ h \end{array}$$

$$h = k$$

3) Če $g \circ f$ mono, je f mono:

$$\text{Ideja: } f \circ h = f \circ k \quad / g \circ -$$

$$g \circ (f \circ h) = g \circ (f \circ k) \quad \text{asoc.}$$

$$\cancel{(g \circ f) \circ h} = \cancel{(g \circ f) \circ k}$$

$$h = k$$

Izrek: Naj bu $f: A \rightarrow B$.

$$1) f \text{ mono} \Leftrightarrow f \text{ injektivna}$$

$$2) f \text{ epi} \Leftrightarrow f \text{ surjektivna}$$

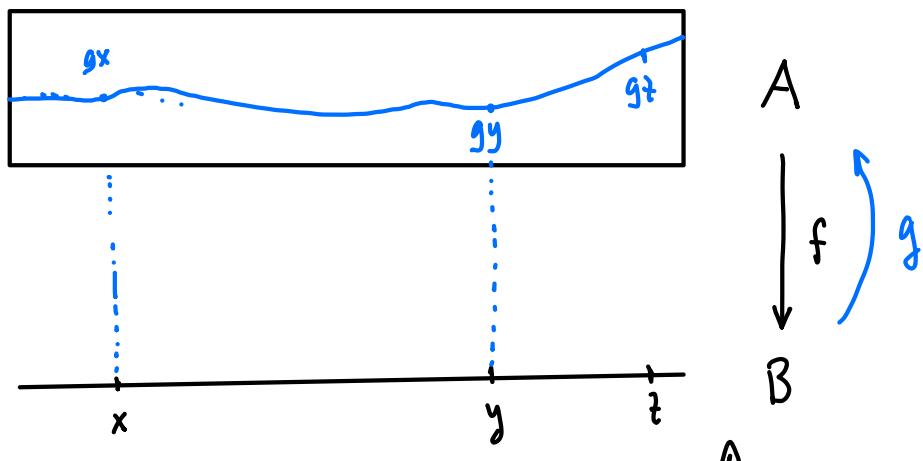
$$3) f \text{ epi \& mono} \Leftrightarrow f \text{ bijektivna} \Leftrightarrow f \text{ izomorfizem}$$

Dohaz: glej zapiske

Retrakcija in prerez:

Def: Če sta $f: A \rightarrow B$ in $g: B \rightarrow A$ takši, da velja $f \circ g = id_B$, potem:

- f je retrakcija ali levi invert g
- g je prerez ali desni invert f .



Izrek: Retrakcija je epi, prerez je mono.

Dokaz: Imamo zgornjo sliko in $f \circ g = id_B$.

$$f \circ g = id_B$$

izrek ⑦

$$id_B \text{ mono} \Rightarrow f \circ g \text{ mono} \Rightarrow g \text{ mono}$$

$$id_B \text{ epi} \Rightarrow f \circ g \text{ epi} \Rightarrow f \text{ epi}$$

□

Izpeljana množica:

Naj bo $f: A \rightarrow B$. Izpeljana množica je

$$\{ f(x) \mid x \in A \} := \{ y \in B \mid \exists x \in A . f(x) = y \}$$

Izpeljana množica s pogojem:

$$\{ f(x) \mid x \in A \wedge \varphi(x) \} := \{ y \in B \mid \exists x \in A . \varphi(x) \wedge f(x) = y \}$$

Običajno pišemo  kut

$$\{ f(x) \mid x \in A \wedge \varphi(x) \}$$

Priimeri:

$$\{ n^2 \mid n \in \mathbb{N} \} = \{ 0, 1, 4, 9, 16, 25, \dots \}$$

$$\{ x^2 + 1 \mid x \in \mathbb{R} \} = [1, \infty)$$

$$\{ n^2 \mid n \in \mathbb{N} \wedge n \text{ prstvilo} \} = \{ 4, 9, 25, 49, 121, \dots \}$$

Slike in pravlike:

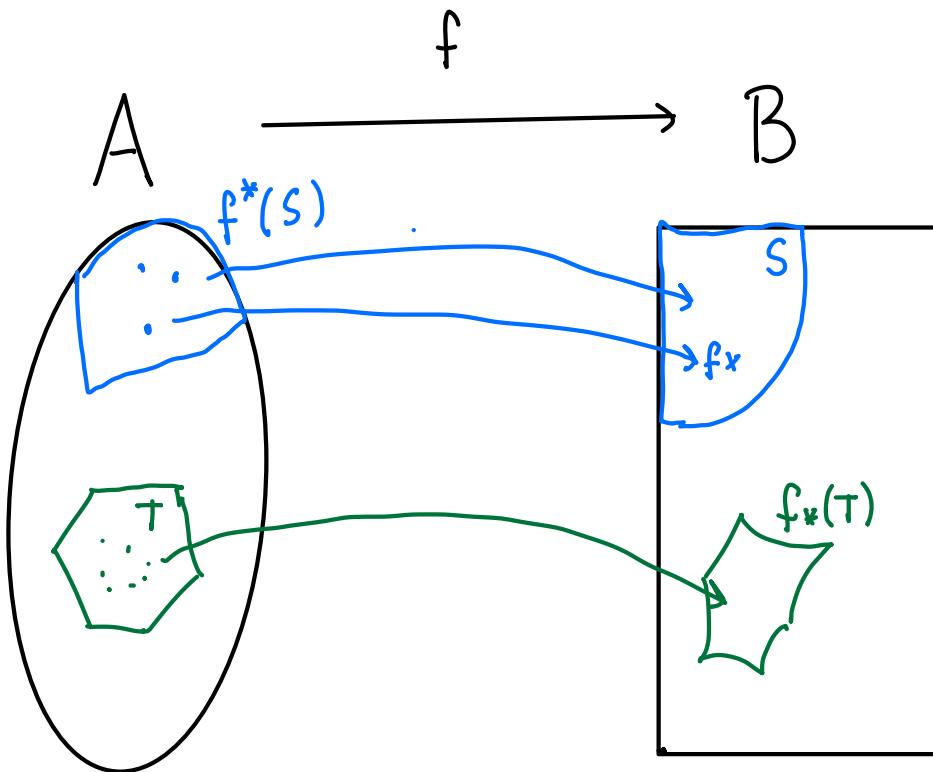
Def: Naj bo $f: A \rightarrow B$.

- pravlika $S \subseteq B$ je $f^*(S) := \{ x \in A \mid f(x) \in S \} \subseteq A$
- slika $T \subseteq A$ je $f_*(T) := \{ y \in B \mid \exists x \in T . f(x) = y \} \subseteq B$
 $= \{ f(x) \mid x \in T \}$

Običajni zapis:

$$\begin{array}{lll} f^*(S) & \text{piše} & f^{-1}(S) \\ f_*(S) & \text{piše} & f(S) \end{array} \left. \begin{array}{c} f^{-1}(S) \\ f(S) \end{array} \right\} \begin{array}{l} \text{zmešuju} \\ f^{-1} \text{ NE POMEVI INVERT } f! \end{array}$$

$f(S)$?!
 $S \subseteq A$ in ne $S \not\subseteq A$.



Def: Zaloga vrednosti $f: A \rightarrow B$ je $f_*(A)$.

Kaj pa $f^*(B)$?

$$f^*(B) = \{x \in A \mid f(x) \in B\} = A$$

$$f^*(\emptyset) = \emptyset$$

$$f_*(A) = B \Leftrightarrow f \text{ surjeftivna} \checkmark$$

Pomini:

$$f^*(T \cup V) = f^*(T) \cup f^*(V)$$

$\cap \qquad \qquad \qquad \cap$

$$f_*(S \cup U) = f_*(S) \cup f_*(U)$$

$$f_*(S \cap U) \subseteq \underline{f_*(S) \cap f_*(U)} \quad \left[\begin{array}{l} \cong \text{velja, i.e.} \\ f \text{ injectivna} \end{array} \right]$$

