

Konstrukcije množic

KARTEZIČNI PRODUKT

A, B množici

$A \times B$ kartezični produkt

Elementi $A \times B$ so urejeni pari

(x, y) če $x \in A$ in $y \in B$

Projekciji

$$\pi_1: A \times B \rightarrow A$$

$$\bullet \pi_1(x, y) = x$$

$$\pi_2: A \times B \rightarrow B$$

$$\bullet \pi_2(x, y) = y$$

$$\text{Velja za } u \in A \times B: \bullet u = (\pi_1 u, \pi_2 u)$$

Primer: $A = \{*, \square\}$ $B = \{1, 2, 3\}$

$$A \times B = \{(*, 1), (*, 2), (*, 3), (\square, 1), (\square, 2), (\square, 3)\}$$

$$\mathbb{N} \times \mathbb{N} = \{(0, 0), (0, 1), (0, 2), (0, 3), \dots, \\ (1, 0), (1, 1), (1, 2), \dots, \\ (2, 0), (2, 1), \dots, \\ \dots\}$$

Primer: $A \times B \cong B \times A$

Izomorfizem: $f: A \times B \rightarrow B \times A$

$$f(x, y) = (y, x)$$

$$g: B \times A \rightarrow A \times B$$

$$g(v, u) = (u, v)$$

Preverimo ali je g inverz f :

$$f(g(v, u)) = f(u, v) = (v, u) \quad \checkmark$$

$$g(f(x, y)) = g(y, x) = (x, y) \quad \checkmark$$

Primer: $(A \times B) \times C \cong A \times (B \times C)$
 $((a, b), c) \quad (a, (b, c))$

vajah ali doma

EKSPONENTNA MNOŽICA

A in B množici

B^A eksponent A in B

elementi B^A : funkcije $A \rightarrow B$.

Primeri: $A = \{1, 2, 3\}$ $B = \{\square, \Delta\}$

$$B^A = \left\{ \begin{array}{c|c} 1 & \square \\ \hline 2 & \square \\ \hline 3 & \square \end{array}, \begin{array}{c|c} 1 & \square \\ \hline 2 & \square \\ \hline 3 & \Delta \end{array}, \begin{array}{c|c} 1 & \square \\ \hline 2 & \Delta \\ \hline 3 & \square \end{array}, \dots, \begin{array}{c|c} 1 & \Delta \\ \hline 2 & \Delta \\ \hline 3 & \Delta \end{array} \right\} \quad \text{osem funkcij}$$

↑
napisi doma manjšoče

$B^A \cong A^B$? V splošnem ne drži.

$A^1 \cong A$ ker imamo izomorfizem: $1 = \{*\}$

$$f: A^1 \rightarrow A$$

$$f(g) = g(*)$$

$$h: A \rightarrow A^1$$

$$x \mapsto 1+x$$

$$y \mapsto y+2-1$$

$$h(a) = (* \mapsto a)$$

Preverimo: $f(h(a)) = h(a)(*) = a$

$$h(f(g)) \stackrel{?}{=} g$$

$$g \in A^1 \quad g: 1 \rightarrow A$$

$$f(g) \in A$$

$$h(f(g)) \in A^1$$

$$h(f(g))(*) = f(g) = g(*) \quad \checkmark$$

FUNKCIJSKI PREDPISI

$$\mathbb{R} \rightarrow \mathbb{R}$$

- $f(x) = x^2 + 7x - 3$

- $x \mapsto x^2 + 7x - 3$

- $\lambda x. x^2 + 7x - 3$

Haskell: $\backslash x \rightarrow x ** 2 + 7 * x - 3$

Python: $\text{lambda } x: x ** 2 + 7 * x - 3$

Primer: $B \rightarrow B^A$
 $y \mapsto (x \mapsto y)$

$$\lambda y. (\lambda x. y)$$

Se enkrat: $B \xrightarrow{f} B^A$
 $f(y) = K_y$

Pomožna definicija:
če je $y \in B$, definiramo
 $K_y: A \rightarrow B$
 $K_y(x) = y$

Izrek: Naj bodo A, B, C množice:

- $A \times \emptyset = \emptyset$
- $A \times 1 \cong A$
- $A \times B \cong B \times A$
- $A \times (B \times C) \cong (A \times B) \times C$
- $A^1 \cong A$
- $A^\emptyset \cong 1$
- $A^{B \times C} \cong (A^B)^C$
- $1^A \cong 1$
- $\emptyset^A \cong \emptyset$ če je A neprazna
- $\emptyset^\emptyset \cong 1$

Dohati: vaje in kasneje

Karakteristične funkcije:

Naj bo $2 = \{0, 1\}$

Izrek: $\mathcal{P}(A) \cong 2^A$

Dohat: Iščevo $\chi: \mathcal{P}(A) \rightarrow 2^A$ izomorfizem.

$$0 = \emptyset$$

$$1 = \{0\}$$

$$2 = \{0, 1\}$$

$$3 = \{0, 1, 2\}$$

$$3 = \{\emptyset, \{\emptyset\}, \{\emptyset, \{\emptyset\}\}$$

Vzemi

$\chi_S \in 2^A$ definiran s predpisom

$$\chi_S(a) = \begin{cases} 0 & \text{če } a \notin S \\ 1 & \text{če } a \in S \end{cases}$$

$$\chi : \mathcal{P}(A) \rightarrow 2^A \quad S \in \mathcal{P}(A)$$

$$\chi_S : A \rightarrow 2 \quad \text{lahko bi pisali } \chi(S)$$

$$\chi_S(a) \in 2, \quad a \in A \quad \text{lahko bi pisali } \chi(S)(a)$$

Trdimo, da je $f : 2^A \rightarrow \mathcal{P}(A)$, definiran s predpisom

$$f(g) = \{a \in A \mid g(a) = 1\} = g^*(\{1\}).$$

Preverimo da sta χ in f inverza:

1) Naj bo $S \in \mathcal{P}(A)$: preverimo $f(\chi_S) = S$.

$$x \in f(\chi_S) \Leftrightarrow$$

$$x \in \{a \in A \mid \chi_S(a) = 1\} \Leftrightarrow$$

$$x \in A \wedge \chi_S(x) = 1 \Leftrightarrow$$

$$x \in A \wedge x \in S \Leftrightarrow$$

$$x \in A \cap S = S$$

2) Naj bo $g \in 2^A$, preverimo $\chi_{f(g)} = g$.

Naj bo $x \in A$:

$$\chi_{f(g)}(x) = \begin{cases} 0 & x \notin f(g) \\ 1 & x \in f(g) \end{cases}$$

$$= \begin{cases} 0 & \neg(x \in A \text{ in } g(x) = 1) \\ 1 & x \in A \text{ in } g(x) = 1 \end{cases}$$

$$\begin{aligned}
&= \begin{cases} 0 & \neg(g(x)=1) \\ 1 & g(x)=1 \end{cases} \\
&= \begin{cases} 0 & g(x)=0 \\ 1 & g(x)=1 \end{cases} \\
&= g(x) \quad \checkmark \quad \blacksquare
\end{aligned}$$

χ_S se imenuje **karacteristična funkcija** podmnožice S .

VSOTA MNOŽIC

(KOPRODUKT, DISJUNKTNA UNIJA)

A in B množici

$A + B$ vsota (pišemo tudi $A \perp B$ in $A \uplus B$)

$$A + B = (\{0\} \times A) \cup (\{1\} \times B)$$

element $A+B$ je bodisi $(0, a)$ za $a \in A$

bodisi $(1, b)$ za $b \in B$

Primer: $A = \{\square, \Delta, \diamond\}$ $B = \{\square, \diamond\}$

$$A + B = \{(0, \square), (0, \Delta), (0, \diamond), (1, \square), (1, \diamond)\}$$

Inkluziji: $\iota_0: A \rightarrow A+B$ $\iota_1: B \rightarrow A+B$
 $\iota_0(x) = (0, x)$ $\iota_1(y) = (1, y)$

Vsak element $A+B$ je bodisi $\iota_0(a)$ za $a \in A$
 bodisi $\iota_1(b)$ za $b \in B$.

Kako definiramo funkcijo $f: A+B \rightarrow C$?

Obraunavamo primere, oziroma po kosih:

$$f(l_0(a)) = g(a) \quad \text{za neki } g: A \rightarrow C$$

$$f(l_1(b)) = h(b) \quad \text{za neki } h: B \rightarrow C$$

Vaje: $A+B \cong B+A$

$$A+\emptyset \cong A$$

$$A+(B+C) \cong (A+B)+C$$

$$A^{B+C} \cong A^B \times A^C$$

$$(A+B) \times C \cong (A \times C) + (B \times C)$$