

Ukazni programski jezik

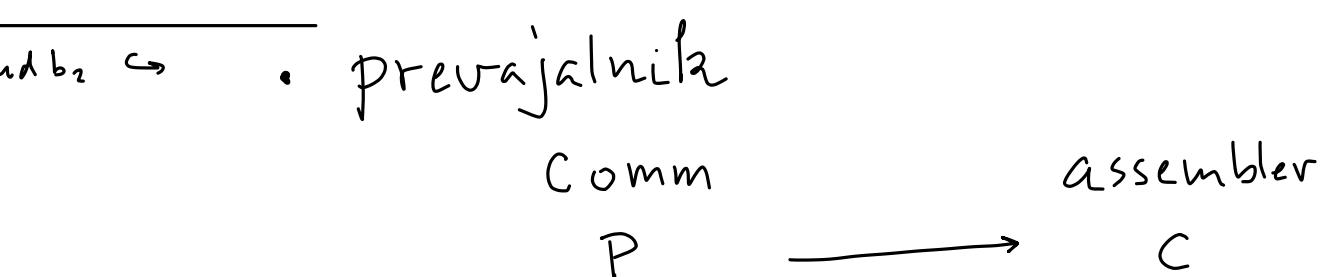
Sintaksa

```

while <bod> do
    <ukaz>; <ukaz>; <ukaz>
done
    <ukaz>
        <ukaz>

```

Operacijska semantika



- potrebujemo opis delovanja, ki je neodvisen od prevajalnika & arhitekture
- Operacijska semantika

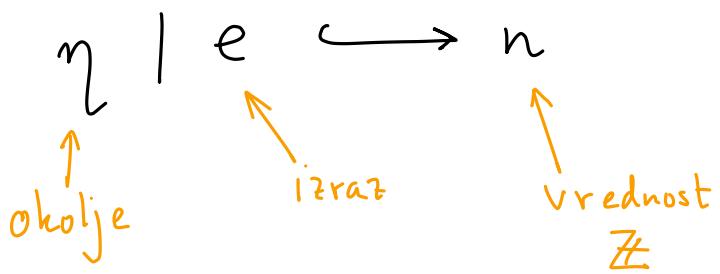
```

def f(n):
    return n*n+3

```

$$\begin{aligned}
 f: \mathbb{Z} &\rightarrow \mathbb{Z} \\
 n &\mapsto n^2 + 3
 \end{aligned}$$

$$f(7) = 52$$

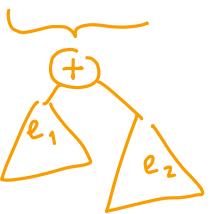


$$\frac{\eta(x) = 10}{\eta \mid x \hookrightarrow 10}$$

$$\eta \mid \underset{3}{x} \overset{\otimes}{\Xi} \hookrightarrow 21$$

$$\eta \mid \underset{x}{\underset{3}{+}} \underset{7}{\Xi} \hookrightarrow 31$$

$10 + 21$



$$\eta \mid b \hookrightarrow p$$

$$\eta \mid \text{not } b \hookrightarrow \neg p$$

$$\eta \mid e \hookrightarrow n$$

$$\eta \mid \neg e \hookrightarrow \neg n$$

b_1 and b_2

$$\frac{\eta \mid b_1 \hookrightarrow \text{false} \quad \eta \mid b_2 \hookrightarrow \text{false}}{\eta \mid b_1 \text{ and } b_2 \hookrightarrow \text{false}}$$

$$\frac{\eta \mid b_1 \hookrightarrow \text{false} \quad \eta \mid b_2 \hookrightarrow \text{true}}{\eta \mid b_1 \text{ and } b_2 \hookrightarrow \text{false}}$$

- "short-circuit"
- izračunaj $b_1 \rightarrow \text{false} \rightarrow$ odgovor: false
 - $\rightarrow \text{true} \rightarrow$ računaj b_2
 - izračunaj b_1 in izračunaj b_2
 - uporabi tabelo

b_1	b_2	$b_1 \text{ and } b_2$
⊥	⊥	⊥
⊥	⊤	⊥
⊤	⊥	⊥
⊤	⊤	⊤

$$\frac{\eta \models b_1 \hookrightarrow \text{true} \quad \eta \models b_2 \hookrightarrow \text{false}}{\eta \models b_1 \text{ and } b_2 \hookrightarrow \text{false}}$$

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Ukazi:

$$\eta \models e \mapsto e'$$

$$(\eta, C) \longmapsto (\eta', C') \text{ en korak}$$

V okoljku η uhat c naredi en korak in novo okolje je η' ter preostaneck uhat je uhat c'

$$(\eta, C) \longmapsto \eta'$$

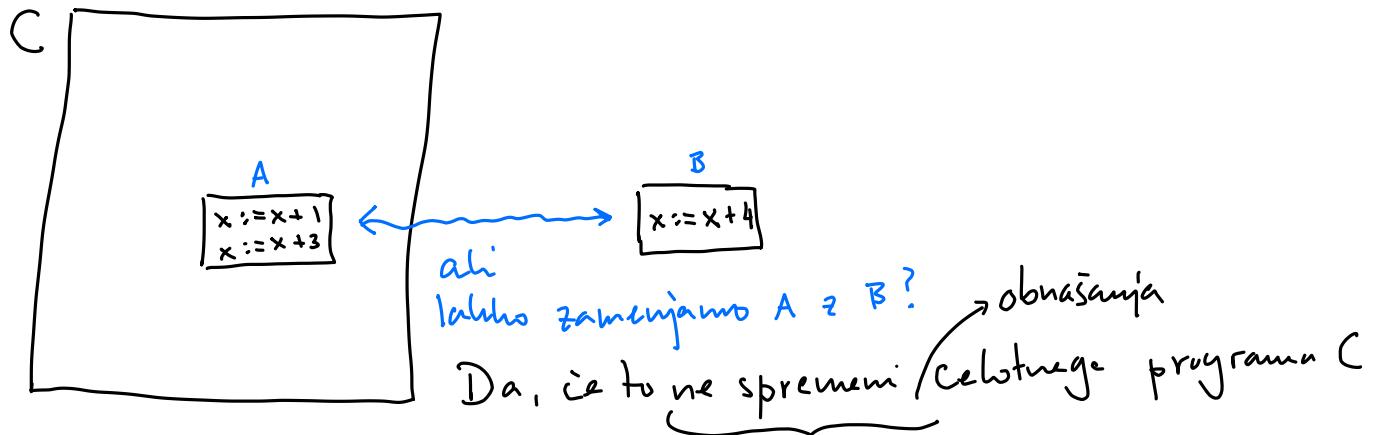
V okoljku η uhat c v enem koraku konča izvajanje in novo okolje je η'

$$\frac{\begin{array}{c} \textcircled{1} \\ (\eta, C_1) \mapsto \eta' \end{array} \quad \begin{array}{c} \textcircled{2} \\ (\eta', C_2) \mapsto (\eta'', C_2') \end{array}}{(\eta, C_1; C_2) \mapsto (\eta'', C_2')}$$

DVA KORAKA!

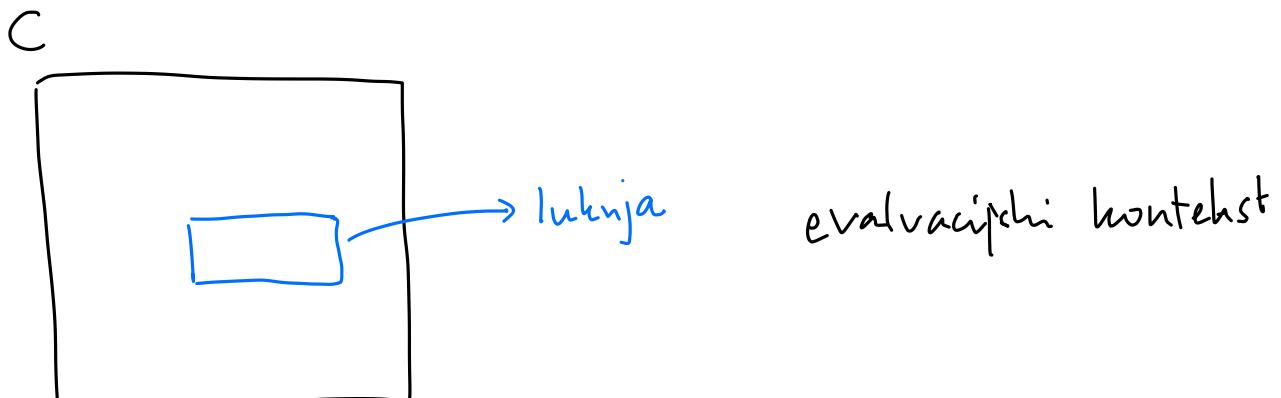
$$\begin{array}{ccc} \text{while } b \text{ do} & C ; & \\ C & \xrightarrow[b=\text{true}]{} & \text{while } b \text{ do} \\ \text{done} & & C \\ & & \text{done} \end{array}$$

Ekvivalenca programov



C bo v obeh primerih izračunal isti rezultat

(ali pa v obeh primerih se zaključi)



for (A ; P ; B) {
 C
}

A;
while (P) {
 C;
 B
}

if p then A else B end

if (p) {
A
}
else {
B
}

if p then A else skip end

if (p) { A }

new x := e in C

nova lokalna spremenljivka
x z začetno vrednostjo e,
veljavna v okazu C

{ int x = e;
C
}