

Recap : Type theory

type / space	A
points	$t : A$
dependent type fibration	$x : A \vdash B(x)$ type $B : A \rightarrow U$ $B(x)$ \vdash universe of types \vdots $\therefore_x A$
unit	$* : 1$
empty	\emptyset
product	$\prod_{x:A} B(x)$ $\lambda(x:A). t(x)$ $x \mapsto t(x)$ A map taking $x:A$ to some point $t(x) : B(x)$
	space of sections $B(x) \rightarrow t(x)$ $\xrightarrow{x} A$
function type	$f : \prod_{x:A} B(x)$ $a : A$ $\overline{fa : B(a)}$ $A \rightarrow B := \prod_{x:A} B$ or B^A

sum $\sum(x:A) B(x)$
 total space $(a, b) : \sum(x:A) B(x)$
 where $a : A$,
 $b : B(a)$

$$\pi_1 : \sum(x:A) B(x) \rightarrow A$$

$$\pi_2 : \prod(u : \sum(x:A) B(x)) . B(\pi_1 u)$$

Natural numbers \mathbb{N}

- $\mathbb{N} : \mathcal{U}$
- Constructors

1) $0_{\mathbb{N}} : \mathbb{N}$

2) if $n : \mathbb{N}$ then $s_n : \mathbb{N}$

3) Induction principle :

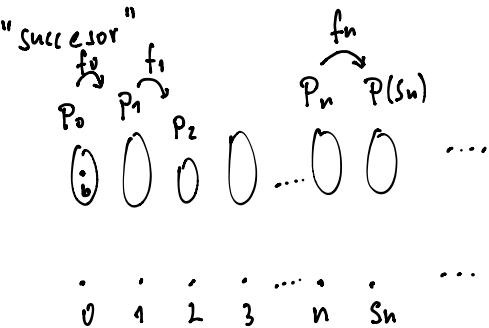
We are given

• $P : \mathbb{N} \rightarrow \mathcal{U}$

• $b : P(0)$

• $f : \prod(n : \mathbb{N}) . P(n) \rightarrow P(s_n)$

• $k : \mathbb{N}$



Have: $\text{ind}_{\mathbb{N}}(P, b, f, k) : P(k)$

Such that: $\text{ind}_{\mathbb{N}}(P, b, f, 0) \equiv_{P(0)} b$

$\text{ind}_{\mathbb{N}}(P, b, f, s_m) \equiv_{P(s_m)} f m (\text{ind}_{\mathbb{N}}(P, b, f, m))$

Recursion principle : $A : \mathbf{U}$

$$P(n) := \begin{array}{c} A \\ A \rightarrow A \\ \vdots \\ b : A \end{array} \quad \begin{array}{c} A & A & A \\ \circlearrowleft & \circlearrowleft & \circlearrowleft \\ f : \prod_{n:N} . A \rightarrow A & = & \dots \\ \downarrow & \downarrow & \downarrow \\ \mathbb{N} \rightarrow A \rightarrow A & & \begin{matrix} 0 & 1 & 2 & 3 & \dots \end{matrix} \end{array}$$

Get : $\text{ind}_{\mathbb{N}}(\lambda_. A, b, f, k) =: \text{rec}_{\mathbb{N}}(A, b, f, k) : A$

Identity Types

In some Quillen model :

$$\begin{array}{ccc} A & \xrightarrow{r_A} & P_A \\ \downarrow \Delta_A & \nearrow \langle \varepsilon_0, \varepsilon_1 \rangle & \text{fibration} \\ A \times A & & \end{array}$$

trivial cofibration

$P_A = A^{[0,1]}$
 $r_A(x) = \lambda t. x$
 $= \text{const}_x$

$\varepsilon_0(r) = r^{(0)}$
 $\varepsilon_1(r) = r^{(1)}$

Consider:

$$\begin{array}{ccc} A & \xrightarrow{d} & B \\ r_A \downarrow & \nearrow \beta = & \downarrow p \text{ fibration} \\ P_A & \xrightarrow{\text{id}} & P_A \\ & \searrow & \swarrow \langle \varepsilon_0, \varepsilon_1 \rangle \\ & A \times A & \end{array}$$

fiber of P_A at $(a, b) \in A \times A$:

$\{r : A^{[0,1]} \mid r(0) = a, r(1) = b\}$

Identity type:

- formation rule: given $A : \mathcal{U}$, $a : A$, $b : A$ there is

$$\text{Id}_A(a, b) : \mathcal{U}$$

Other notations: $a =_A b$ (not to confuse
 $a = b$ with $a \equiv_A b$)

- constructor: given $a : A$ there is

$$\text{refl}_A(a) : \text{Id}_A(a, a)$$

Other notations:

$$\text{idpath}_A(a) : a =_A a$$

- path induction:

Given:

$$\rightarrow B : \prod(x : A)(y : A)(p : x =_A y). \mathcal{U}$$

For all $x : A$, $y : A$, $p : x =_A y$, have $B(x, y, p) : \mathcal{U}$.

$$\rightarrow d : \prod(z : A). B(z, z, \text{refl}_A(z))$$

$$\rightarrow a : A, b : A, q : a =_A b$$

Have: $J(B, d, a, b, q) : B(a, b, q)$

$$J(B, d, c, c, \text{refl}_A(c)) \equiv_{B(c, c, \text{refl}_c)} d(c) \quad \text{for } c : A$$

Path induction for humans:

- have $B(x, y, p)$ depends on $x, y : A$ and $p : x =_A y$
NB: x, y and p must be unrestricted / arbitrary.
- Want point of $\prod(x, y : A) (p : x =_A y) B(x, y, p)$
- Sufficient to give a point of $\prod(z : A) B(z, z, \text{refl}_A(z))$.

Example: Paths can be inverted

$$\text{inv} : \prod(x, y : A). x =_A y \rightarrow y =_A x$$

$$\text{Solution: } \prod(x : A). \prod(y : A). \prod(p : x =_A y). y =_A x$$

$$\text{Take } B(x, y, p) := y =_A x$$

$$\text{inv} : \prod(x, y : A) (p : x =_A y) . B(x, y, p)$$

By path induction enough to give

$$d : \prod(z : A). B(z, z, \text{refl}_A(z))$$

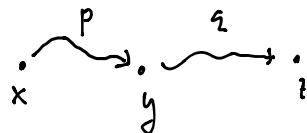
$$\prod(z : A), z =_A z$$

$$\text{Take } d = \lambda(z : A). \text{refl}_A(z).$$

$$\text{Get } \text{inv}(x, y, p) : y =_A x \quad \text{Write } \bar{p}' := \text{inv}(x, y, p)$$
$$\begin{array}{c} \uparrow \\ x =_A y \end{array}$$

$$\text{Know: } (\text{refl}_A(z))^{-1} \equiv \text{refl}_A(z).$$

Exercise: Paths can be composed



$$\text{TT}(x, y : A)(p : x = y)(z : A)(q : y = z) \cdot x = z \\ \underbrace{\hspace{10em}}_{B(x, y, p)}$$

Path induction:

$$\text{TT}(v : A) B(v, v, \text{refl}(v)) \\ d : \text{TT}(v : A) \cdot \text{TT}(z : A)(q : v = z) \cdot N = z$$

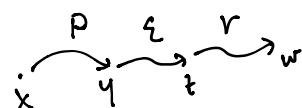
$$\text{Take } d = \lambda_{N : A} \cdot \lambda_{z : A} \cdot \lambda q : v = z \cdot q$$

$$\text{Write } p * q := J(B, d, x, y, p)(z)(q) \quad \text{for} \\ p : x = y, q : y = z$$

$$\text{Know: } \underset{A}{\text{refl}(x)} * q \equiv_{x=z} q \quad \text{for } q : x = z$$

Path composition is associative:

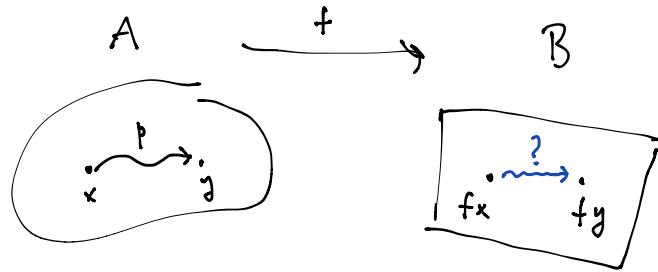
$$p * (q * r) = \underset{x = A w}{(p * q) * r}$$



$$\text{Conclusion: } \sum(x : A) \cdot \sum(y : A) \cdot x =_A y$$

$$(x, y, p) \quad \text{where } \begin{array}{c} x : A \\ y : A \end{array}$$

$$p : x =_A y$$



$$\prod (x:A)(y:A)(p:x=_A y). \underbrace{f x =_B f y}_{B(x,y,p)}$$

Path induction

$$\prod (z:A). f z =_B f z$$

$$\lambda z. \text{refl}_B(f z)$$

Notation: $\text{ap}_f(p) : f x =_B f y$
 $f(p)$

Know: $f(\text{refl}_z) \equiv \text{refl}_B(f z)$

Exercises:

$$f(p^{-1}) = f(p)^{-1}$$

$$f y =_B f x$$

$$f(p * g) = f(p) * f(g)$$

