

# Univalence

① Path spaces - how to "compute" them?

type  $A$ ,  $x, y : A$

What does  $x =_A y$  look like?

Example  $\prod(x, y : 1). (x =_1 y) \simeq 1$ .  $* : 1$

Construct  $f : \prod(x, y : 1). x =_1 y \rightarrow 1$

$g :$   $\leftarrow$

$f x y p := *$

$g x y u : x =_1 y$  induction on  $x : 1$  and  $y : 1$

$g * * u : * =_* *$   
 $:= \text{refl} *$

Verify:  $f x y (g x y u) = u$  ? for all  $x, y, u : 1$

Induction on  $u : 1$ :

$f x y (g x y *) = *$

$*$  by definition of  $f$

Verify:  $g * y (f * y p) = p$  for all  $x, y : 1$ ,  $p : x =_1 y$ ?

Induction on  $p$ :

$$x := ?$$

$$y := ?$$

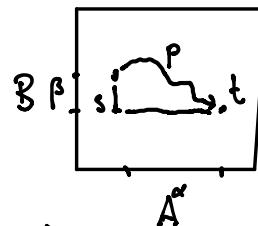
$$p := \text{refl } ?$$

$g ? ? (f ? ? (\text{refl } ?)) = \text{refl } ?$  induction on  $? : 1$

$g * * (f * * (\text{refl } *)) \stackrel{\substack{\uparrow \\ \text{def of } g}}{=} \text{refl } * \checkmark$

□

Example:  $s, t : A \times B$



$$\prod(s, t : A \times B). (s =_{A \times B} t) \simeq (\pi_1 s =_A \pi_1 t) \times (\pi_2 s =_B \pi_2 t)$$

$$p : s = t \mapsto (\pi_1(p), \pi_2(p))$$

$$? \leftarrow (\alpha, \beta)$$

Example:

Recall:

There is a map

$$X \cong Y \xrightarrow{\epsilon} X \simeq Y.$$

If you just need any equivalence  $e : X \cong Y$ , then it is enough to construct  $i : X \cong Y$ , then

$$e := \epsilon(i).$$

Example:  $f, g: A \rightarrow B$

$f =_{A \rightarrow B} g$  what it's like?

Would want to show:

$$(f =_{A \rightarrow B} g) \simeq \text{TT}(x:A). f x =_B g x$$

Example: Say  $A$  has decidable equality when  
 $\neg x = x \rightarrow 0$   
 $\text{TT}(x,y:A). (x =_A y) + \neg(x =_A y)$ .

Theorem (Hedberg): if  $A$  has decidable equality then  $A$  is a set.

② What about the path spaces in a universe?

$A, B : \mathcal{U}$  What is  $A =_{\mathcal{U}} B$  like?

$$\{\text{jaka}\} = \{\text{pope}\}$$

"Equivalent types are equal."

Want an equivalence

$$(A =_{\mathcal{U}} B) \simeq (A \simeq B)$$

"Equality is equivalent to equivalence".

$\text{idtoeq} : \text{TT}(A, B : \mathcal{U}) . A =_{\mathcal{U}} B \rightarrow A \simeq B$

path induction:

$\text{TT}(C : \mathcal{U}) . C \simeq C$

$C \mapsto \text{idtoeq}_C$  the identity equivalence on  $C$

Given  $p : A =_{\mathcal{U}} B$  we get  $\text{idtoeq}_{A,B} p : A \simeq B$ .

Univalence axiom ( $\text{UA}$ ):

$\text{TT}(A, B : \mathcal{U}) . \text{isequiv}(\text{idtoeq}_{A,B})$ .

This amounts to:

- there is a map

$\text{ua}_{A,B} : A \simeq B \rightarrow A =_{\mathcal{U}} B$

- there is a map

$\text{idtoeq}_{A,B} : A =_{\mathcal{U}} B \rightarrow A \simeq B$

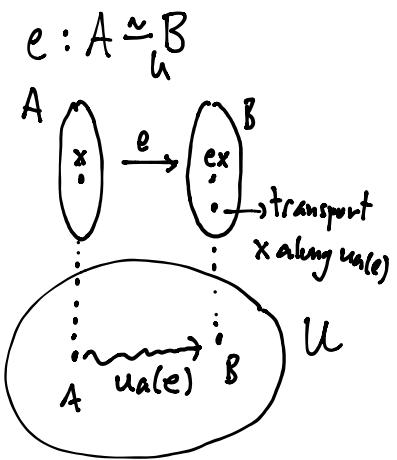
- Inverses:

$$\text{ua}(\text{idtoeq } p) =_{A =_{\mathcal{U}} B} p$$

$$\text{idtoeq}(\text{ua}(e)) =_{A \simeq B} e$$

- Computation rule:

$$\text{transport}^{(\lambda x : \mathcal{U}. X)}_{(u \text{al}(e), x)} =_B e x$$



In general:

$B : A \rightarrow \mathcal{U}$  is univalent when

$$\prod(x, y : A). (x =_A y) \simeq (Bx \simeq By)$$

$$Bx \xrightarrow{\sim} By$$

$$\bigcirc \quad \bigcirc$$

$$\begin{array}{c} \xrightarrow{x} \\ \text{---} \\ y \end{array} \quad A$$

When is  $B : 1 \rightarrow \mathcal{U}$  univalent?

$$\prod(x, y : 1). (x =_1 y) \simeq (Bx \simeq By)$$

$$(=_1, *) \simeq (B* \simeq B*)$$

$$1 \simeq (B* \simeq B*)$$

Plug in  $\text{Id} : \mathcal{U} \rightarrow \mathcal{U}$  to get  $UA$ .

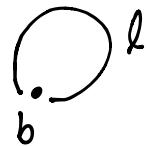
$$\textcircled{3} \quad \pi_1(S^1) = \mathbb{Z}$$

Recall:

- circle  $S^1$

$b: S^1$  base

$\ell: b = b$  loop



- $(\mathbb{Z}, +, 0)$  the additive group of integers

- fundamental group:

$A$ ,  $a: A$  loop space at  $a$ :

$$\Omega(A, a) := (a =_A a)$$

$$\pi_1(A, a) := \|\Omega(A, a)\|_0$$

We show that  $\Omega(S^1, b) = \mathbb{Z}$ , so it follows

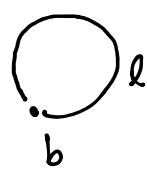
$$\pi_1(S^1, b) = \|\Omega(S^1, b)\|_0 = \|\mathbb{Z}\|_0 = \mathbb{Z}$$

Enough to show  $\Omega(S^1, b) \cong \mathbb{Z}$ . by UA.

$$(b =_{S^1} b) \cong \mathbb{Z}$$

$$f: \mathbb{Z} \rightarrow (b =_{S^1} b)$$

$$g: (b =_{S^1} b) \rightarrow \mathbb{Z}$$



Candidate for  $f$ :

$$f(k) := \begin{cases} \underbrace{\text{loop} \cdot \text{loop} \cdot \dots \cdot \text{loop}}_k & \text{if } k > 0 \\ \text{refl}_b & \text{if } k = 0 \\ \underbrace{\text{loop}^1 \cdot \dots \cdot \text{loop}^1}_{-k} & \text{if } k < 0 \end{cases} =: \text{loop}^k$$

Candidate for  $g$ :



$$g(p) := ? \text{ winding number}$$

$$\text{Verify: } \prod_{b=b} f(g(p)) = p.$$

Problem: Cannot use path-induction on  $p$ .

Define:

$$C : S^1 \rightarrow \mathcal{U}$$

$$D : S^1 \rightarrow \mathcal{U}$$

$$\boxed{\begin{array}{l} S^1 \rightarrow X \\ b \mapsto x \\ l \mapsto p : x =_X x \end{array}}$$

as follows:

- Define  $D$ :  $D(x) := (b =_{S^1} x)$  universal cover

- Define  $C$  by circle induction:

$$C(b) := \mathbb{Z}$$

"deck transformations"

$$C(\text{loop}) : C(b) =_{\mathcal{U}} C(b)$$

$$C(\text{loop}) : \mathbb{Z} =_{\mathcal{U}} \mathbb{Z}$$

...

$$\text{ual}(\text{succ}) : \mathbb{Z} =_u \mathbb{Z} \quad \text{because} \quad \begin{aligned} \text{succ}(k) &= k+1 \\ \text{pred}(k) &= k-1 \end{aligned}$$

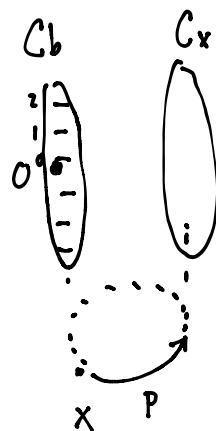
$$\text{succ} : \mathbb{Z} \simeq \mathbb{Z}$$

Plan:  $\prod(x:S'). D_x \simeq C_x$

$$\prod(x:S'), (b =_{S'} x) \simeq C_x$$

In particular when  $x \equiv b$ , we get

$$(b =_{S'} b) \simeq \mathbb{Z}$$



Define:  $e : \prod(x:S'). D_x \rightarrow C_x$

$$e \times p := \text{transport}^C(p, 0)$$

$$\downarrow$$

$$b =_{S'} x$$

Define:  $d : \prod(x:S'). C_x \rightarrow D_x$

by circle induction

$$\begin{aligned} d(b) : C_b &\rightarrow D_b \\ \mathbb{Z} &\rightarrow b=b \end{aligned}$$

$$d(b) := \lambda k : \mathbb{Z}. \text{loop}^k$$

Also need to lift the loop: (find a path above the loop)  
SKIP

Verify:

$$d x (e \times p) = p ? \quad p : b =_{S^1} x$$

for all  $x : S^1$ , all  $p : b = x$ . Path induction on  $p$ :

$$d b (e b (\text{refl } b)) =$$

$$d b (\text{transport}^c(\text{refl } b, 0)) =$$

$$d b 0 = \text{loop}^b = \text{refl } b \quad \checkmark$$

Verify:  $e \times (d x c) = c$  for all  $x : S^1$   
 $c : Cx$

Circle induction: verify for  $x = b$

$$e b (d b c) = c ? \quad c : Cb \equiv \mathbb{Z}$$

for all  $c$

$$e b (d b c) = e b (\text{loop}^c)$$

$$= \text{transport}^c(\text{loop}^c, 0)$$



$$= \text{succ}^c(0)$$

$$= c$$

